Nash Bargaining Solution and Application

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Networks: Technology, Economics, and Social Interactions

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Outline

- Part I: Theory
 - Bargaining Problem
 - Nash Bargaining Solution
- Part II: Application
 - Case: Mobile data offloading
- Conclusion

Bargaining Problem

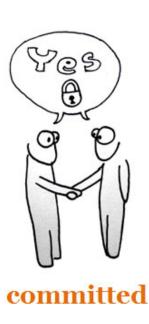
- Bargaining is one of the most common activities in daily life.
 - Price bargaining in an open market;
 - Wage and working time bargaining in a labor market;
 - Score bargaining after an examination;
 - ·



bargaining

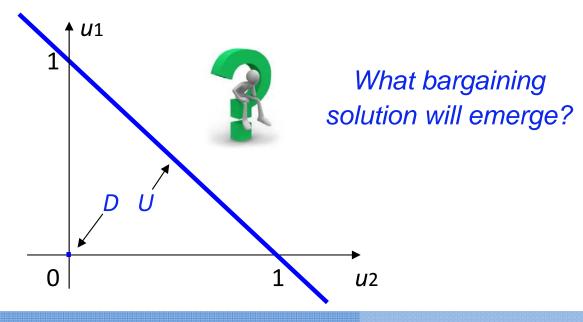
Bargaining Problem

- Formally, bargaining problems represent situations in which:
 - Multiple players with specific objectives search for a mutually agreed outcome (agreement).
 - No agreement may be imposed on any player without his approval, i.e., the disagreement is possible.
 - Players have the possibility of reaching a mutually beneficial agreement.
 - There is a conflict of interest among players about agreements.
- Bargaining solution
 - An agreement or a disagreement



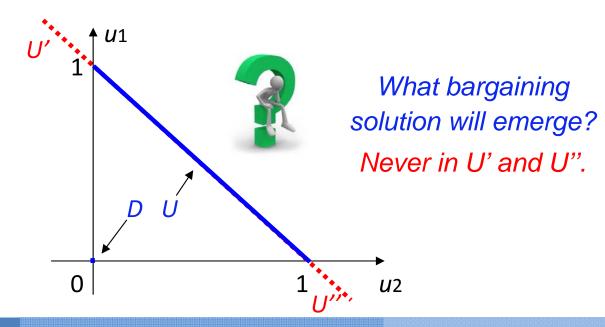
A simple example

- A simple example: player 1 wants to sell a book to player 2.
 - Problem: Players bargain for the price p
 - The objective (payoff) of players: $u_1=p$, $u_2=1-p$
 - The set of feasible agreements: $U = \{(u_1, u_2) \mid u_1 + u_2 = 1\}$
 - The disagreement: $D = (d_1, d_2)$, e.g., D=(0,0)
 - A bargaining solution is an outcome $(v_1, v_2) \subseteq U \cup D$



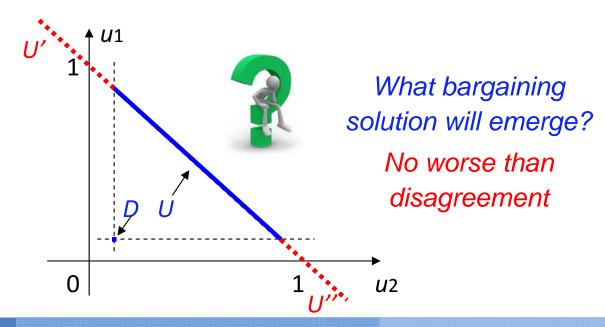
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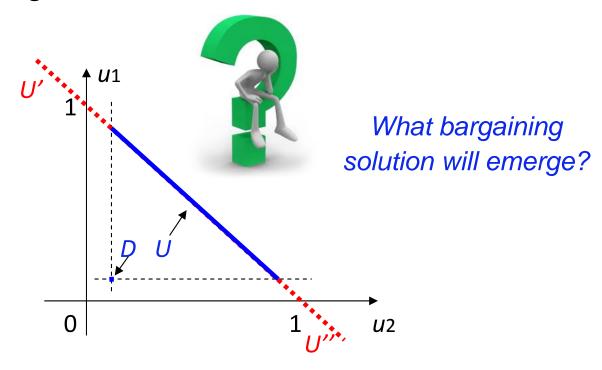


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 - The set of feasible agreements: $U = \{(u_1, u_2) \mid u_1 + u_2 = 1\}$
 - The disagreement: $D = (d_1, d_2)$, e.g., D = (0.1, 0.1)
 - A bargaining solution is an outcome $(v_1, v_2) \subseteq U \cup D$



- Bargaining theory is a theoretic tool used to identify the bargaining solution, given
 - (i) the set of all feasible agreements U
 - (ii) the disagreement D



- Strategic Approach vs Axiomatic Approach
 - Strategic approach: (i) Modeling the bargaining process (i.e., the game form) explicitly, and (ii) Considering the game outcome (i.e., equilibrium) that results from the players' self-enforcing interactions.
 - e.g., Rubinstein Bargaining Model, 1982
 - Axiomatic approach: (i) Abstracting away the details of the process of bargaining, and (ii) Considering only the set of outcomes or agreements that satisfy "reasonable" properties.
 - e.g., Nash Bargaining Model, 1950

- Bargaining solution by strategic approach
 - A simple illustration: Player 1 wants to sell a book to player 2
 - Stage 1: Player 1 proposes a price $p=p_1$, and player 2 accepts or refuses; If accept, bargaining terminates; If not, turn to Stage 2;
 - Stage 2: Player 2 proposes a price $p=p_2$, and player 1 accepts or refuses; If accept, bargaining terminates; If not, turn to Stage 3;
 - Stage 3: Player 2 proposes a price $p=p_3$, and player 2 accepts or refuses; If accept, bargaining terminates; If not, turn to Stage 4;
 -

The bargaining solution is the equilibrium of this game. Example: Rubinstein Bargaining Model, 1982

- Bargaining solution by axiomatic approach
 - A simple illustration: Player 1 wants to sell a book to player 2
 - Axiom 1: Pareto efficiency.
 - Axiom 2: Equal share of payoff gain.
 - Axiom 3: ...
 - ·

The bargaining solution is the solution satisfying all axioms.

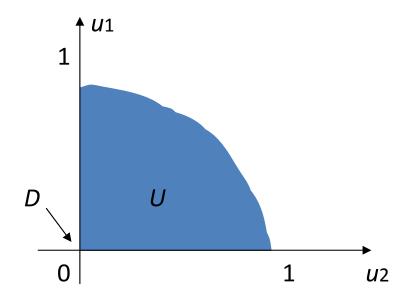
Example: Nash Bargaining Model, 1950 Shapley Bargaining Model, 1976

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- 2-person bargaining problem [Nash J., 1950]
- An axiomatic approach based bargaining solution
- 4 Axioms
 - (1) Pareto Efficiency
 - (2) Symmetry
 - ▶ (3) Invariant to Affine Transformations
 - ► (4) Independence of Irrelevant Alternatives
- Nash bargaining solution is the unique solution that satisfies the above 4 axioms.

2-person Bargaining Problem

- A general 2-person bargaining model
 - The set of bargaining players: $N = \{1,2\}$
 - ► The set of feasible agreements: $U = \{(u_1, u_2) \in a \text{ bounded convex set}\}$
 - The outcome of disagreement: $D = (d_1, d_2)$, e.g., D = (0,0)
 - A Nash bargaining solution is an outcome $(v_1, v_2) \subseteq U \cup \{D\}$ that satisfies the Nash's 4 axioms.



Nash's Axioms

Nash's 4 Axioms

- ► (1) Pareto Efficiency: None of the players can be made better off without making at least one player worse off;
- (2) Symmetry: If the players are indistinguishable, the solution should not discriminate between them;
- ▶ (3) Invariant to Affine Transformations: An affine transformation of the payoff and disagreement point should not alter the outcome of the bargaining process;
- ▶ (4) Independence of Irrelevant Alternatives: If the solution (v_1, v_2) chosen from a feasible set A is an element of a subset $B \subseteq A$, then (v_1, v_2) must be chosen from B.

Thought: Are these axioms reasonable?

Nash bargaining solution (NBS) is the unique solution that satisfies the Nash's 4 axioms.

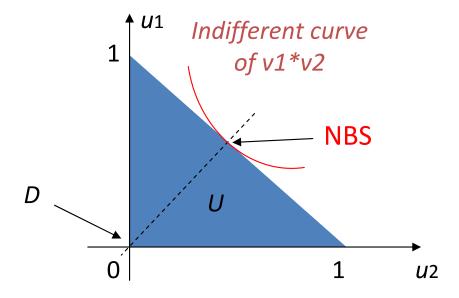
Definition

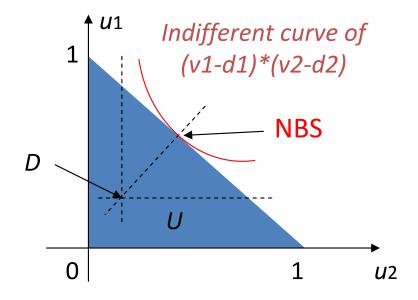
We say that a pair of payoffs (v_1^*, v_2^*) is a Nash bargaining solution if it solves the following optimization problem:

$$\max_{v_1, v_2} \quad (v_1 - d_1)(v_2 - d_2)$$
subject to
$$(v_1, v_2) \in U$$

$$(v_1, v_2) \ge (d_1, d_2)$$

- An illustration of NBS: 2 players split 1 dollar
 - The set of feasible agreements: $U = \{(u_1, u_2) \mid u_1 + u_2 \le 1, u_1, u_2 \ge 0\}$
 - The outcome of disagreement: $D = (d_1, d_2)$





(a) NBS when
$$D=(0,0)$$

 $(v1,v2) = (0.5, 0.5)$

(b) NBS when
$$D=(0.3,0.2)$$

 $(v1,v2) = (0.55, 0.45)$

- Important factors determining a NBS
 - Feasible agreement sets U
 - Disagreement D
 - ► Increase a player's disagreement → higher payoff for the player in Nash bargaining solution.
 - Bargaining power a
 - ► Increase a player's bargaining power → higher payoff for the player in Nash bargaining solution.

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Case: Mobile Data Offloading

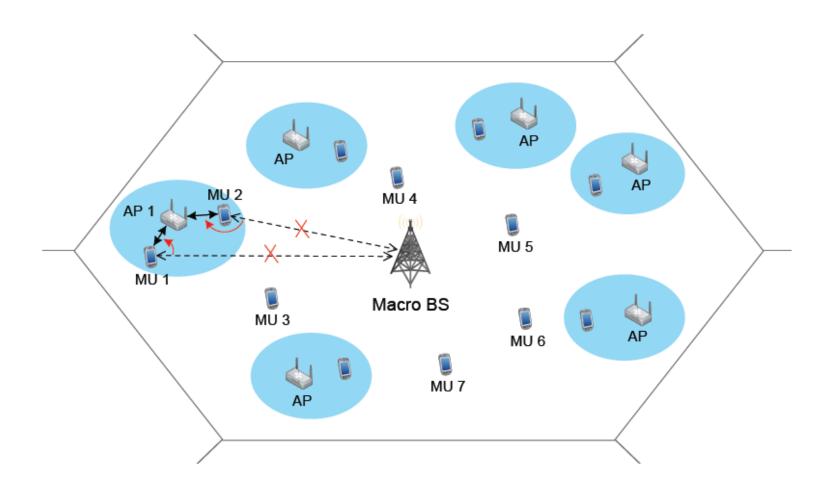
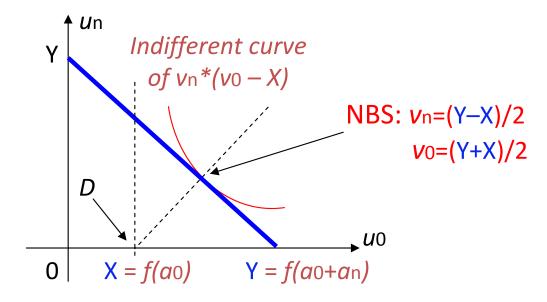


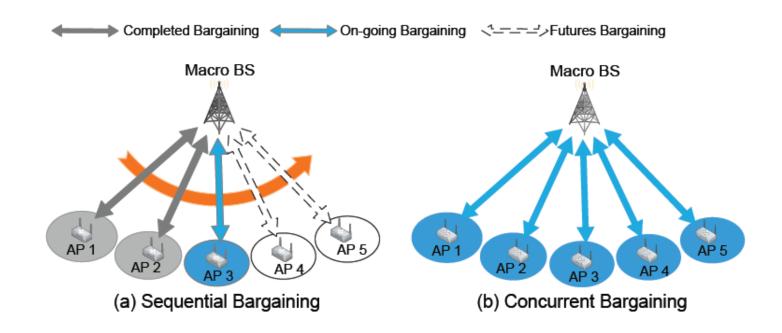
Figure. An illustration of mobile data offloading in a macrocell BS. A Mobile User (MU) can directly access the BS or offload his data (send or receive) through an AP if he is within both coverage areas.

- BS wants to buy spectrum resource from APs
 - Problem: BS and each AP n bargain for the price pn
- Key notations
 - The amount of AP n's resource: an, n=1,2,...,N
 - The price for AP n's resource: p_n , n=1,2,...,N
 - The amount of BS's own resource: ao
 - The welfare function of BS: f(a) (strict increasing)

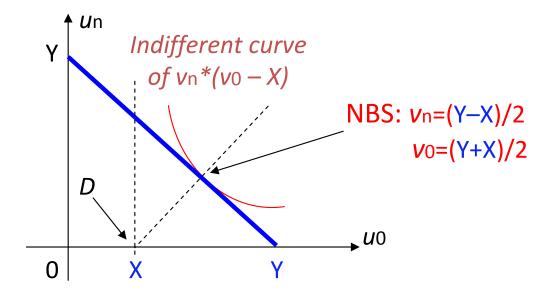
- An illustration of bargaining between AP n and BS
 - 2-person bargaining problem
 - The objective (payoff) of players: $u_n = p_n$, $u_0 = f(a_0 + a_n) p_n$
 - The set of feasible agreements: $U = \{(u_n, u_0) | u_n + u_0 = f(a_0 + a_n)\}$ (blue curve)
 - The disagreement: $D = (0, f(a_0))$



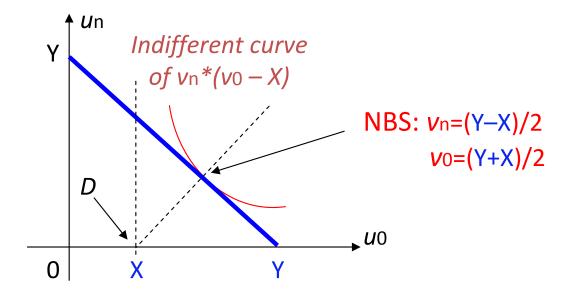
- Totally N two-person bargaining problems
- Bargaining protocol:
 - Sequentially bargaining
 - Concurrently bargaining



- Sequentially Bargaining in step n: BS and AP n
 - 2-person bargaining problem
 - The objective (payoff) of players: $u_n = p_n$, $u_0 = f(a_0 + a_1 + ... + a_n) p_n$
 - The set of feasible agreements: $U = \{(u_n, u_0) \mid u_n + u_0 = f(a_0 + a_1 + ... + a_n)\}$
 - The disagreement: $D = (0, f(a_0+a_1+...+a_{n-1}))$



- Concurrently Bargaining between BS and AP n
 - 2-person bargaining problem
 - The objective (payoff) of players: $u_n = p_n$, $u_0 = f(a_0 + a_1 + ... + a_N) p_n$
 - The set of feasible agreements: $U = \{(u_n, u_0) \mid u_n + u_0 = f(a_0 + a_1 + ... + a_N)\}$
 - The disagreement: $D = (0, f(a_0+a_1+...+a_{n-1}+a_{n+1}+...+a_N))$



- Examples: 10 APs
 - The amount of AP n's resource: $a_n = 1$, n=1,2,...,10
 - ightharpoonup The amount of BS's own resource: $q_0 = 5$
 - ▶ The welfare function of BS: $f(a) = \log(a)$
- NBS in Sequentially Bargaining
 - APs: $u_1 = 0.5*\log(6) 0.5*\log(5)$, $u_2 = 0.5*\log(7) 0.5*\log(6)$, ..., $u_{10} = 0.5*\log(15) 0.5*\log(14)$,
 - **BS**: $u_0 = 0.5*\log(15) + 0.5*\log(5)$
- NBS in Concurrently Bargaining
 - APs: $u_1 = 0.5*\log(15) 0.5*\log(14)$, $u_2 = 0.5*\log(15) 0.5*\log(14)$, ..., $u_{10} = 0.5*\log(15) 0.5*\log(14)$,
 - BS: $u_0 = 5*\log(14) 4*\log(15)$

Conclusion

- We discuss the basic theory of bargaining solution, in particular the Nash bargaining solution.
- We discuss a potential application of Nash bargaining solution in wireless networks: mobile data offloading.

Thank you!