### **Enabling Crowd-sourced Mobile Internet Access**

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Fig. Illustration of Crowd-Sourced Mobile Internet Access: A set of mobile devices form a group/community (via WiFi or Bluetooth), sharing internet connections with each other.



Fig. Illustration of Crowd-Sourced Mobile Internet Access: A device may act as a gateway/host downloading data for others, a relay forwarding data for others, and a client consuming data.

# Key Problems in This Work

- How to achieve an efficient and fair network resource allocation?
  - who (hosts) will download data for whom (clients), and how much?
  - who (relays) will route data from each host to each client, and how much?
- How to encourage the user participation and cooperation?
  - how to compensate the hosts and the relays for their efforts?

### Outline









# **Mobile Data Traffic Explosion**



Fig. Global Mobile Data Traffic, 2012 to 2017 (from Cisco VNI)

Mobile data traffic explosive growth: 66% annual grow rate

Reaching 11.2 exabytes per month by 2017, a 13-fold increase over 2012 or a 46-fold increase over 2010.

### How about the network?



Fig. Historical Increases in Spectral Efficiency (from Femtoforum)

- Network capacity slow growth: less than 29% annual grow rate
  - Available spectrum band growth: 8% per year
  - Cell site increase: 7% per year
  - Spectrum efficiency growth: less than 12% per year from 2007 to 2013

 $108\% \cdot 107\% \cdot 112\% = 129\%$ 

### How to deal with the data explosion?

Network capacity growth vs Data traffic growth

29% vs 66%

### • Traditional approach: Network Expansion

- Upgrading access technology (e.g., WCDMA  $\rightarrow$  LTE  $\rightarrow$  LTE-A)
- Acquiring new spectrum license (e.g., TV white space)
- Building more pico/micro/macro cell sites
- However, all of these methods are costly and time-consuming.

### How to deal with the data explosion?

- Data usage is highly imbalanced, e.g., time-, location-, and network-dependent.
- One promising approach: Network Outsourcing
  - (A) Outsource network to network (N2N)
    - ★ Example: Mobile data offloading/onloading
  - ► (B) Outsource network to mobile user (N2MU)
    - ★ Example: User-provided networking



• In this work, we focus on the scenario B — N2MU.

### A General Model of N2MU

- A general model: A set of mobile users form a group/community, sharing internet connections with each other.
  - ----- Crowd-Sourced Mobile Internet Access



- Key Features
  - (i) User-provided networking
    - \* Mobile users can access internet through the hosting of other users.
  - (ii) Multi-hop accessing
    - \* Mobile users can access internet through the relay of multiple devices.
  - (iii) Access bonding
    - \* Mobile users can access internet through multiple access links.



- Roles of mobile users
  - Gateway (Host): Downloading data for others
  - Relay: Forwarding data for others
  - Client: Consuming data



- Real Cases
  - Open Garden (http://opengarden.com)



M-87 (http://www.m-87.com/)



# Key Problems in This Work

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- How to encourage the user participation and cooperation?
  - how to compensate the hosts and the relays for their efforts?
- Our purpose: Design a mechanism to address both the incentive issue and the efficiency and fairness issues, i.e.,
  - Encourage the user participation and cooperation;
  - Achieve an efficient and fair network resource allocation.









### 4 Conclusion

- A set of I mobile users:  $\mathcal{I} = \{1, 2, ..., I\}$ 
  - Each user may act as a gateway/host downloading data for others, a relay forwarding data for others, and a client consuming data.
  - Users are heterogeneous in terms of their data requirements, internet connections, and energy budgets.



 $\bullet$  Parameters related to each user  $i \in \mathcal{I}$ 



- $c_i, c_{ij}, c_{ji, j \in \mathcal{I}}$ : link capacity;
- $e_i, e_{ij}^s, e_{ij, i \in \mathcal{I}}^r$ : unit energy consumption;
- $\triangleright$  p<sub>i</sub>: unit internet connection cost.

• Variables related to user  $i \in \mathcal{I}$  as a client.



▶  $y_{j, j \in I}^{(i)}$ : the data downloaded via other hosts for client i; ▶  $y^{(i)} = y_1^{(i)} + ... + y_I^{(i)}$ : the total data consumed by client i; ▶  $U_i(y^{(i)})$ : the utility function of client i.

• Variables related to user  $i \in \mathcal{I}$  as a host (gateway).



- $y_{i,i\in\mathcal{I}}^{(j)}$ : the data downloaded via host i for other clients;
- $y_i = y_i^{(1)} + ... + y_i^{(I)}$ : the total data downloaded via host i;
- $\blacktriangleright \ e_i \cdot y_i:$  the total energy consumption for downloading data;
- $p_i \cdot y_i$ : the total payment for downloading data;
- Downloading capacity constraint:  $y_i \leq c_i$ .

• Variables related to user  $i \in \mathcal{I}$  as a relay.



- ▶  $x_{ij, n \in I}^{(n)}$ : the data relayed from user i to user j, for client n;
- $e_{ij}^r \cdot \sum_n x_{ji}^{(n)}$ : the total energy consumption for receiving data from user j;
- $e_{ij}^{s} \cdot \sum_{n} x_{ij}^{(n)}$ : the total energy consumption for sending data to user j.

• Variables related to user  $i \in \mathcal{I}$  as a relay.



- ▶ Relaying capacity constraint:  $\sum_{n} x_{ij}^{(n)} \le c_{ij}$ ,  $\sum_{n} x_{ji}^{(n)} \le c_{ji}$ ▶ Flow balance constraint:  $\sum_{i} x_{ji}^{(n)} + y_{i}^{(n)} = \sum_{i} x_{ij}^{(n)}$ ,  $n \in \mathcal{I}$
- - \* For each relay i, its total received data (of client n) equals its total sent data (of client n).

• Payoff of each user  $i \in \mathcal{I}$ :

$$J_i(\textbf{x}_i, \textbf{y}_i) = U_i - P_i - E_i$$

• 
$$\mathbf{x}_i = \{x_{ij}^{(n)}\}_{j,n\in\mathcal{I}}$$
: Relaying matrix;

- $\mathbf{y}_i = \{y_i^{(n)}\}_{n \in \mathcal{I}}$ : Downloading matrix;
- U<sub>i</sub>: Utility of user i;
- P<sub>i</sub>: Total payment of user i (for internet access);
- E<sub>i</sub>: Total energy consumption of user i;

### Outline

1 Background





### 4 Conclusion

### **Purpose of This Study**

#### **Objective**

Our purpose is to design a mechanism to address the incentive, efficiency, and fairness issues in crowd-sourced mobile internet access, including

- Encouraging the user participation and cooperation;
- Achieving an efficient and fair network resource allocation.

### **Incentive Issue**

- How to encourage the user participation and cooperation?
  - Users may not want to participate in the crowd-sourced system;
    - ★ For example, those without the current demand;
  - Users may not want to download or relay data for others, especially those helpless to them.
    - \* For example, user i may not want to download data for user 4.



### **Incentive** Issue

### • Our Solution: Virtual Currency

Key idea: User pays certain virtual currency to those who send data to him (i.e., I give you money, you give me data).



z<sup>(n)</sup><sub>ji</sub>: the virtual price that user i pays j for receiving data (of client n);
∑<sub>n</sub> z<sup>(n)</sup><sub>ji</sub> ⋅ x<sup>(n)</sup><sub>ji</sub>: the total virtual money that user i pays j for receiving data (of all clients) from j;

### **Modified Payoff with Virtual Currency**

• Modified payoff of each user  $i \in \mathcal{I}$ :

$$J_i(\textbf{x}_i,\textbf{y}_i,\textbf{z}_i) = U_i - P_i - E_i + V_i$$

i;

• 
$$\mathbf{z}_i = \{z_{ij}^{(n)}\}_{j,n\in\mathcal{I}}$$
: Virtual price matrix;  
•  $V_i$ : Total virtual currency evaluation of user

### **Efficiency and Fairness Issues**

- How to achieve an efficient and fair network resource allocation?
  - Efficiency: The aggregate payoff of all users is maximised or close to the maximum;
  - Fairness: Every user achieves a satisfactory payoff;
- Our Solution: Nash Bargaining Solution
  - ► Key idea: Users bargain for the network resource allocation and the virtual currency transfer based on the Nash bargaining framework.

# Nash Bargaining Solution

### Nash Bargaining Problem (NBP)

 $\begin{array}{ll} \displaystyle \max_{\textbf{x}_i,\textbf{y}_i,\textbf{z}_i,\forall i} & \Pi_{i\in\mathcal{I}} \big(J_i-J_i^0\big) \\ \\ \text{s.t.,} & (a) \ J_i \geq J_i^0 & (J_i^0: \text{disagreement point}) \\ & (b) \ \textit{Capacity constraints}; \\ & (c) \ \textit{Flow balance constraint}; \\ & (d) \ \textit{Virtual current budget constraint}. \end{array}$ 

#### • The NBP problem has a unique optimal solution.

# Nash Bargaining Implementation

#### • Centralized Implementation

 A central control node collects all the required network information, and computes the Nash bargaining solution.

#### • Decentralized Implementation

- Iterative updating: Users update their individual decisions sequentially and repeatedly, where in each updating period,
  - one user first updates its decision based on a local objective function and signals from its one-hop neighbors,
  - \* and then reports the necessary signals to its one-hop neighbors.

### Simulation

- An example with 6 nodes
  - Blue Bar: Downloading/relaying data;
  - Red Bar: Consuming data;



Left: Independent Operation.



0.337 0.239

0.410

Operation

.247

5

0.046

### Outline

1 Background

2 Model





### Conclusion

- We study the crowd-sourced mobile internet access system, in particular, we answer
  - How to achieve an efficient and fair network resource allocation?
  - How to encourage the user participation and cooperation?
- We propose a Nash bargaining solution with virtual currency, which addresses the incentive, efficiency, and fairness issues.

### **Future Extension**

- Interference is a big problem !
  - Intra-node interference (RF constraint)
    - ★ E.g., links 3, 4, and 5.
  - Inter-node interference (Transmission collision)
    - ★ E.g., links 1 and 2.



# Thank You



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